## RSA Cryptography

Alice wants to receive a secure one-word message from Bob

- Alice Chooses prime numbers p,q whose product is greater than 26.
- (in practice, these would be very large numbers, and their product would be huge.)
- Ann computes $\mathrm{n}=\mathrm{pq}$
- Alice chooses a positive integer e that is relatively prime to (p-1)(q-1)
- Alice computes the an integer $d$ that is a positive multiplicative inverse of $e, \bmod (p-1)(q-1)$
- The numbers n,e are called the Public Key. Alice sends Bob the $n$ and the e, the Public Key. (Alice does not send Bob the values of $\mathrm{p}, \mathrm{q}, \mathrm{d}$.)

Bob has a word consisting consisting of $k$ letters chosen from the set $\{\mathrm{a}, \mathrm{b}, \ldots, \mathrm{z}\}$. The letters are denoted $\mathcal{L}_{1}, \mathcal{L}_{2}, \ldots, \mathcal{L}_{k}$ (The word would be written with the letters side-by-side with no commas, $\mathcal{L}_{1} \mathcal{L}_{2} \ldots \mathcal{L}_{k}$ )

- Bob receives the Public Key $n, e$ from Alice.
- Bob repeats the following steps for each letter $\mathcal{L}_{j}$ in his word, for $j=1,2, \ldots k$
- He converts the letter $\mathcal{L}_{j}$ to a number in the range 1-26, called $M_{j}$
- Then he computes $\left(M_{j}\right)^{e} \bmod n$. The result is denoted $C_{j}$. So

$$
C_{j}=\left(M_{j}\right)^{e} \bmod n
$$

- Bob sends Alice the list of numbers $C_{1}, C_{2}, \ldots, C_{k}$

Alice

- Alice receives the list of numbers $C_{1}, C_{2}, \ldots, C_{k}$ from Bob
- Alice repeats the following steps for each number $C_{j}$ in the list, for $j=1,2, \ldots k$
- She computes $\left(C_{j}\right)^{d} \bmod n$. The result is $M_{j}$.That is,

$$
M_{j}=\left(C_{j}\right)^{d} \bmod n
$$

- She converts converts the number $M_{j}$ to letter $\mathcal{L}_{j}$
- The result is a list of letters $\mathcal{L}_{1}, \mathcal{L}_{2}, \ldots, \mathcal{L}_{k}$
- The resulting word is $\mathcal{L}_{1} \mathcal{L}_{2} \ldots \mathcal{L}_{k}$


## Observations:

- Alice does not send p,q,or d. She only sends Bob n and e.
- Without knowing the value of d, one cannot decrypt Bob's message.
- And without knowing the values of p, q, one cannot find d.
- One could guess values of $p, q$ by factoring $n$. But in practice, $n$ is a very large number, and so factoring n is not feasable in a reasonable time scale.

